CS 201 midterm review

10/7/18

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What topics should we go over?

55 responses



Agenda

- 1. How to get better at UNIX
- 2. cons vs. append
- 3. Deep recursion (review lambda functions)
- 4. Boolean functions
- 5. Turing machines (review let)
- 6. Time permitting: more questions!

Please interrupt at any time with questions or suggestions :)

UNIX

How can I get better at UNIX?

1) UNIX tutorial on the Zoo! ssh into the Zoo; then in your home folder, type the following command: python3 /c/cs201/www/unixtutorial.py

2) Practice typing commands on the Zoo

General tips:

• Be familiar with the *output* of each command (important in context of the transcript!)

Any specific UNIX questions?

cons vs. append

First, a note on pairs vs. lists...

From the Racket guide:

- "A pair combines exactly two values."
- "A list is recursively defined: it is either the constant null, or it is a pair whose second value is a list."

```
(list v ...) \rightarrow list? procedure
v : any/c
```

Returns a newly allocated list containing the vs as its elements.

```
> (list 1 2 3 4)
'(1 2 3 4)
```

cons

(cons a d) → pair?
 a : any/c
 d : any/c

Returns a newly allocated pair whose first element is *a* and second element is *d*.

```
△ (cons 1 2) returns the pair `(1 . 2)
This pair is not a list because the cdr is not null!
```

```
If you want to construct the list `(1 2), use
(cons 1 (cons 2 `())) or
(cons 1 `(2)) or
(cons 1 (list 2))
```

cons continued

In general, here is how I use/conceptualize cons to return a list (not merely a pair):

- Let a be any data type; let b be a list
- (cons a b) inserts a as the first element of the list b

Examples:

```
> (cons 1 '())
'(1)
> (cons 'apple '(banana cranberry))
'(apple banana cranberry)
> (cons '(1 2) (list 3 4))
'((1 2) 3 4)
```

Key takeaway:

<u>The second argument for cons should almost</u> <u>always be a list</u> (unless you want to return a pair); the first argument can be whatever you want and will be inserted as the first element of the list supplied

append

(append lst ...) → list?
 lst : list?
(append lst ... v) → any/c
 lst : list?
 v : any/c

When given all list arguments, the result is a list that contains all of the elements of the given lists in order. The last argument is used directly in the tail of the result.

procedure

The last argument need not be a list, in which case the result is an "improper list."

In general, here is how I use/conceptualize append to return a list:

- Let a and b be lists
- (append a b) essentially merges the lists a and b

```
> (append '(1) '())
'(1)
> (append '(apple) '(banana cranberry))
'(apple banana cranberry)
> (append '((1 2)) '(3 4))
'((1 2) 3 4)
```

Key takeaway:

<u>Generally speaking, all arguments for</u> <u>append should be lists</u> (unless you want to return an improper list)

```
> (cons 1 '())
'(1)
> (cons 'apple '(banana cranberry))
'(apple banana cranberry)
> (cons '(1 2) (list 3 4))
'((1 2) 3 4)
> (append '((1 2)) '(3 4))
'((1 2) 3 4)
```

Questions?

Practice with cons vs. append

- 1. (cons 1 2)
- 2. (cons 1 '())
- 3. (cons 1 '(2))
- 4. (cons '(1) 2)
- 5. (cons '(1) '(2))
- 6. (append '(1) '(2))
- 7. (append `((1)) `((2)))
- 8. Define my-lst to be `(hello ``hi" #t 7).
 - a. How would you use cons to get the list `(1 hello "hi" #t 7)?
 - b. How would you use append to get the list `(1 hello "hi" #t 7)?

Solutions

- 1. (cons 1 2) => `(1 . 2) 2. (cons 1 `()) => `(1) 3. (cons 1 `(2)) => `(1 2) 4. (cons `(1) 2) => `((1) . 2) 5. (cons `(1) `(2)) => `((1) 2) 6. (append `(1) `(2)) => `(1 2) 7. (append `((1)) `((2))) => `((1) (2)) 8. Define my-lst to be `(hello ``hi" #t 7).
 - a. How would you use cons to get the list `(1 hello ``hi" #t 7)?=>
 (cons 1 my-lst)
 - b. How would you use append to get the list `(1 hello ``hi" #t 7)?
 => (append `(1) my-lst)

Deep recursion

Practice with deep recursion

Write a procedure

```
(count-if pred tree)
```

which returns the number of leaves of the tree that satisfy the given predicate pred

Examples

```
(count-if odd? '(1 2 3)) => 2
(count-if even? '(1 2 3)) => 1
(count-if integer? '(1 (2 (3)))) => 3
(count-if string? '()) => 0
(count-if even? '((((((8 8 8))))))) => 3
(count-if (lambda (x) (> x 5)) '(((((((9 9 9))))))) => 3
```

Quick review: what is a lambda function?

A lambda expression creates a function. In the simplest case, a lambda expression has the form

(lambda (arg-id ...)
 body ...+)

```
Example: > ((lambda (x y) (+ x y)) 17 4)
21
> (define (add-x-to-y x y) (+ x y))
> (add-x-to-y 17 4)
21
```

Back to deep recursion...

```
Write a procedure
```

```
(count-if pred tree)
```

which returns the number of leaves of the tree that satisfy the given predicate pred

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Back to deep recursion...

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(count-if (lambda (x) (> x 5)) '((((((9 9 9))))))) => 3
```

Sample solution

```
(define (count-if pred tree)
  (cond [(empty? tree) 0]
      [(list? (first tree)) (+ (count-if pred (first tree)) (count-if pred (rest tree)))]
      [(pred (first tree)) (+ 1 (count-if pred (rest tree)))]
      [else (count-if pred (rest tree))]))
```

Sample solution

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Questions? Comments? Do you agree or disagree?

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```

Questions? Comments? Do you agree or disagree?

Time to draw a tree of recursive calls

(count-if (lambda (x) (and (negative? x) (odd? x))) '(((-13) 4) (-57) 6))

Sample tree of recursive calls



Boolean functions

Key ideas

- Truth tables
- Operations:
 - and *
 - or +
 - not `
- Sum of products algorithm

Truth table example – write an expression for f(x,y,z)

x	У	Z	f(x,y,z)
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

Truth table example – write an expression for f(x,y,z)How Lapproach truth tables:

1. Find all the true values in the output column

x	У	Z	f(x,y,z)
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	
1	0	0	0
1	0	1	
1	1	0	1
1	1	1	

Truth table example – write an expression for f(x,y,z)

	f(x,y,z)	Z	У	x
	0	0	0	0
	0	1	0	0
-	0	0	1	0
x'*y*z		1	1	0
•	0	0	0	1
x*y'*z		1	0	1
x*y*z′		0	1	1
x*y*z	1	1	1	1

- 1. Find all the true values in the output column
- 2. Write Boolean expressions for the corresponding rows

Truth table example – write an expression for f(x,y,z)

	f(x,y,z)	Z	У	x
	0	0	0	0
	0	1	0	0
	0	0	1	0
x′*y*z		1	1	0
	0	0	0	1
x*y'*z		1	0	1
x*y*z'		0	1	1
x*y*z		1	1	1

- 1. Find all the true values in the output column
- 2. Write Boolean expressions for the corresponding rows
- Add these expressions together to get your final sum of products: (x'*y*z)+(x*y'*z)+(x*y*z')+(x*y*z)

Hooray! You're done! You've written a Boolean expression for f(x,y,z) using the sum-of-products algorithm

> Equivalently: (x*y)+(y*z)+(x*z) Can you see why this is true?

Practice with truth tables – write an expression for f(x,y,z)

x	У	Z	f(x,y,z)
0	0	0	1
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	1
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Practice with truth tables – write an expression for f(x,y,z)

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0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	0

Solution: (x'*y'*z')+(x'*y*z')+(x*y'*z')+(x*y'*z)

Turing machines

First, a quick review of let - what is it good for?

A let form binds a set of identifiers, each to the result of some expression, for use in the let body:

(let ([*id expr*] ...) *body* ...+)

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Example from my hw3.rkt:

What's the advantage here?

Think/pair/share – what is this Turing machine doing?

```
(define tm-mystery
 (list (ins 'q1 0 'q1 0 'R)
        (ins 'q1 1 'q1 1 'R)
        (ins 'q1 'b 'q2 'b 'L)
        (ins 'q2 0 'q3 1 'L)
        (ins 'q2 1 'q4 0 'L)
        (ins 'q3 0 'q3 1 'L)
        (ins 'q3 1 'q4 0 'L)
        (ins 'q3 'b 'q5 'b 'R)
        (ins 'q4 0 'q4 0'L)
        (ins 'q4 1 'q4 1 'L)
        (ins 'q4 'b 'q5 'b 'R)))
```

(Assume inputs will be >= 1 in binary)

Solution

Subtract 1 from the n-bit input; the output is n-bit difference

Examples: 1 => 0 10 => 01 11 => 10

11 -> 10

100 => 011

Questions?